

Simplicial completion of (pre)model categories and strictification (Joint work in progress with Chaitanya Leena Subramaniam)

Simon Henry

University of Ottawa

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Simplicial completion is a process to turn a general model category \mathcal{C} into a simplicial one. This is done by considering the category $s\mathcal{C}$ of simplicial objects in \mathcal{C} .

Theorem (Dugger)

If C is a left proper combinatorial Quillen model category, then the category sC of simplicial objects of C has a simplicial model structure such that:

- *The constant object functor $C \rightarrow sC$ is a left Quillen equivalence.*
- *Any left Quillen functor $C \rightarrow D$ with D a simplicial model category can be factored uniquely into a simplicial left Quillen functor $C \rightarrow sC \rightarrow D$.*

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First, sC is a simplicially enriched, tensored and cotensored category:

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$$\text{Hom}^{\widehat{\Delta}}(X, Y)_k = \int_{[n] \in \Delta} \text{Hom}(X([n]), Y([n]))^{\Delta([n],[k])} \quad Y^A = \int_{[n] \in \Delta} Y([n])^{A([n])}$$

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We will identify \mathcal{C} with the subcategory of $s\mathcal{C}$ of constant objects.

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- Generating cofibrations:

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$$\Lambda^k[n] \otimes B \coprod_{\Lambda^k[n] \otimes A} \Delta[n] \otimes A \rightarrow \Delta[n] \otimes B$$

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The model structure on sC is constructed as a left Bousfield localization of the Reedy model structure whose local objects are the Reedy fibrant homotopically constant diagram $\Delta^{\text{op}} \rightarrow C$. It is easy to see that this is Quillen equivalent to C .

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Theorem (Dugger + Barwick, Batanin - White)

If C is a *combinatorial Quillen model category*, then the category sC of simplicial objects of C has a *simplicial left semi-model structure* such that:

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Let C be a κ -combinatorial premodel category. Let $\text{Cof}_\kappa C$ be the category of κ -presentable cofibrant objects.

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- 1 F sends the initial object to $\{*\}$.
- 2 F sends pushout along cofibrations to pullback.
- 3 F sends κ -small transfinite composition of cofibration to limits.
- 4 F sends anodyne cofibrations to weak equivalences.

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Term Equality axioms:

$$x, y : \text{Ob}; f : \text{Hom}(x, y) \vdash f \circ_{x,x,y} \text{Id}_x = f \quad x, y : \text{Ob}; f : \text{Hom}(x, y) \vdash \text{Id}_y \circ_{x,y,y} f = f$$

$$w, x, y, z : \text{Ob}; f : \text{Hom}(w, x); g : \text{Hom}(x, y); k : \text{Hom}(y, z) \vdash k \circ (g \circ f) = (k \circ g) \circ f$$

Given such a theory T , we can form a syntactic category C_T of T , whose objects are “context” and whose morphism are definable “function” between such context up to provable equality

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$$\Delta = (a, b : \text{Ob}; v : \text{Hom}(a, b))$$

$$\Gamma = (x, y, z : \text{Ob}; g, f : \text{Hom}(x, y); h : \text{Hom}(y, x); k : \text{Hom}(z, x))$$

are contexts.

$$\begin{array}{ccc} \Gamma & \rightarrow & \Delta \\ (x, y, z, g, f, h, k) & \mapsto & (z, y, f \circ k) \end{array}$$

is a context morphism.

This category C_T has a special class of maps that I will call “display maps” that corresponds to just forgetting the last variable from the context. For example:

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These maps are closed under pullbacks.

Theorem (Cartmell)

The category of models of T is equivalent to the category of functor $C_T \rightarrow \text{Set}$ preserving the terminal object and pullback along display maps.

In particular, the Yoneda embedding define a fully faithful functor $Y : C_T^{\text{op}} \rightarrow \text{Mod}(T) \subset \text{Fun}(C_T, \text{Set})$.

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The category $\text{Cof}_\omega(\text{Mod}(T))$ of finitely presentable cofibrant objects in $\text{Mod}(T)$ is equivalent to (the Cauchy completion) C_T^{op} .

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- 3 Equivalently, fibrant objects are the simplicial models $C_T \rightarrow s\text{Set}$ which takes values in Kan complexes and send each display map to a Kan fibration. (So all pullbacks appearing in the definition of pullbacks are pullback of fibrations)
- 4 Weak equivalence between fibrant objects are the levelwise weak equivalences.

We can apply our main theorem to this premodel category $\text{Mod}(T)$. We obtain:

Theorem (H. - L. S.)

The localization of the model category $s\text{Mod}(T)$ of strict simplicial model of T at weak equivalence is equivalent to the ∞ -category of weak models of T .

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The localization of the model category $s\text{Mod}(T)$ of strict simplicial model of T at weak equivalence is equivalent to the ∞ -category of weak models of T .

Where a “weak model” is a (pseudo) functor $C_T \rightarrow \text{Spaces}$ that preserve the terminal objects sends pullback of display map to (homotopy) pullbacks.

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The ∞ -category of simplicial abelian groups is equivalent to the ∞ -category of (weak) model of the Lawvere theory of Abelian groups. Note that these are not the same as the (weak) models of the commutative operads (i.e. E_∞ -algebras).

Example

Taking T the theory of category mentioned before $s\text{Mod}(T)$ is a Model structure on simplicial categories whose fibrant objects are simplicial categories C where the simplicial set of objects C_0 is a Kan complex and the map $C_1 \rightarrow C_0 \times C_0$ is a Kan fibration.

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C_T^{op} is the full subcategory of Cat of categories that are free on graphs (so it contains Δ). And the functors we consider are the one that satisfies an analogue of the Segal conditions (their value on any graph is recovered as a limit of $F(\Delta[0])$ and $F(\Delta[1])$ over all vertices and edge of the graph).

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$$R : \widehat{D} \rightleftarrows C : N_D$$

To solve this sort of problem, we have a more precise version of the theorem, where we can input a proposed notion of weak models.

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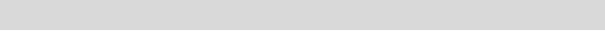
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Example

Take $C = \text{Mod}(\text{Cat})$ and $D = \Delta$ (seen as a dense full subcategory of categories). And the localization $s\widehat{\Delta}_{\text{Proj}}^{\text{Segal}}$.



We also assume that C is κ -combinatorial and that it has a set of generating cofibration in D .

Theorem (H. - L. S.)

In the situation above, Then given S a set of maps in $s\widehat{D}$, the functors R and N_D induces a Quillen equivalence between sC and the localization $s\widehat{D}_{Prof}^{S-Loc}$ if and only if:

- *All maps in S are sent to isomorphism by the derived functor $\overline{R} : \widehat{D} \rightarrow Ho(sC)$.
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This is not trivial, but is well known.

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A version of this in the special case of "pasting diagrams" (instead of all positive polygraphs) has been established last year by Tim Campion.

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In particular, there is a Bousfield localization of sC (i.e. strict simplicial n -categories) which models weak (∞, n) -categories. The case $n = 1$ is exactly Horel's model structure on simplicial categories.

All of this works more generally when C is an extensive category and Θ is replaced by any set of connected objects.